

Lightning Network

Formal verification of a payment protocol

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03/04/2024



Context



Bitcoin's lack of scalability → off-chain protocols



Protocol not proven → attacks



Our goal : prove the security of a fix

Core idea



Lock coins on the chain



Exchange this money off-chain



Use the chain to cash in

Opening a channel

Published on-chain



Funding transaction	
<u>Input</u> Alice: 5 Bob: 5	<u>Output</u> A&B: 10

Saved off-chain



Closing transaction	
<u>Input</u> A&B: 10	<u>Output</u> Alice: 5 Bob: 5

Updating a channel



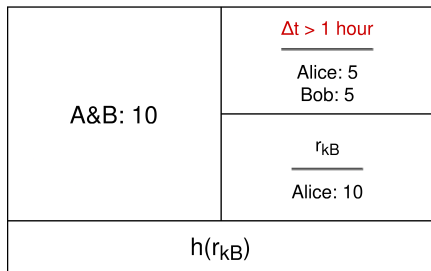
Old transaction	
A&B: 10	Alice: 5 Bob: 5



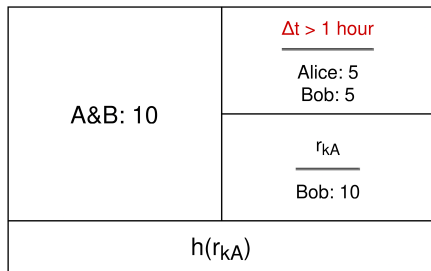
New transaction	
A&B: 10	Alice: 4 Bob: 6

Revocation mechanism

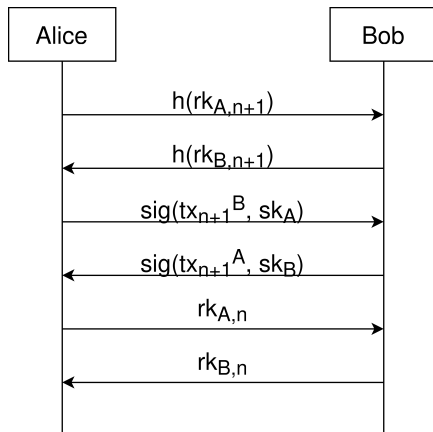
Signed by Alice, held by Bob



Signed by Bob, held by Alice

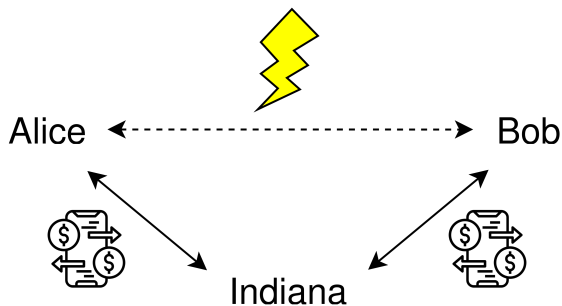


Revocation mechanism



Lightning Network

What if Alice and Bob don't share a channel ?



Hashed Timelock Contract

Signed by Alice, held by Indiana

A&I: 10	$\Delta t > 1 \text{ hour}$ <hr/> Alice: 5 Indiana: 5
	$s; t < 18\text{h}00$ <hr/> Alice: 4 Indiana: 6
	rk_I <hr/> Alice: 10
$h(s), h(rk_I)$	

Signed by Indiana, held by Bob

I&B: 17	$\Delta t > 1 \text{ hour}$ <hr/> Indiana: 12 Bob: 5
	$s; t < 17\text{h}00$ <hr/> Indiana: 11 Bob: 6
	rk_B <hr/> Indiana: 17
$h(s), h(rk_B)$	

Hashed Timelock Contract

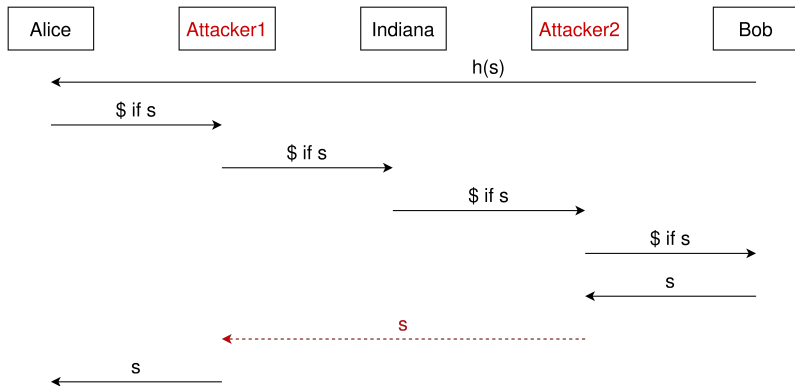
Signed by Alice, held by Indiana

A&I: 10	$\Delta t > 1 \text{ hour}$ <hr/> Alice: 5 Indiana: 5
	$s; t < 18\text{h}00$ <hr/> Alice: 4 - fee Indiana: 6 + fee
	rk_I <hr/> Alice: 10
$h(s), h(rk_I)$	

Signed by Indiana, held by Bob

I&B: 17	$\Delta t > 1 \text{ hour}$ <hr/> Indiana: 12 Bob: 5
	$s; t < 17\text{h}00$ <hr/> Indiana: 11 Bob: 6
	rk_B <hr/> Indiana: 17
$h(s), h(rk_B)$	

Wormhole attack



Desired properties

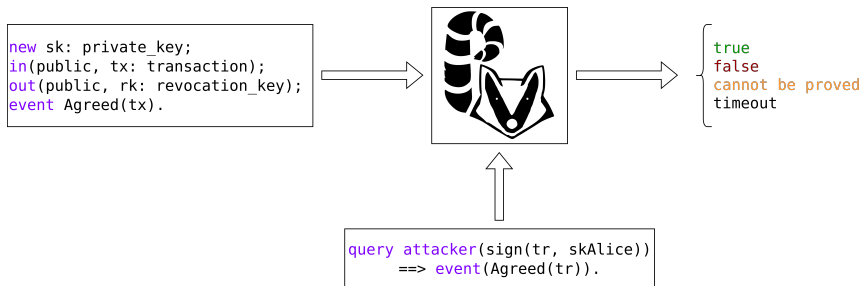


Honest participants cannot lose money

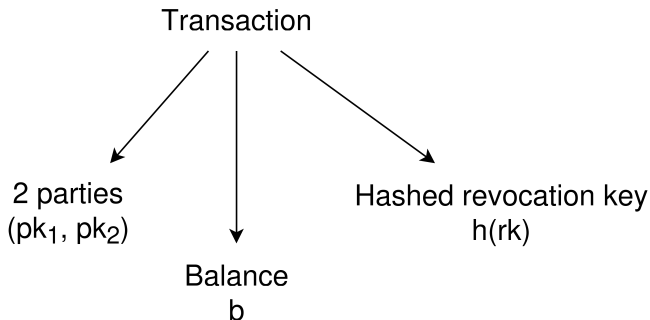


Honest participants get their fees

The Proverif tool

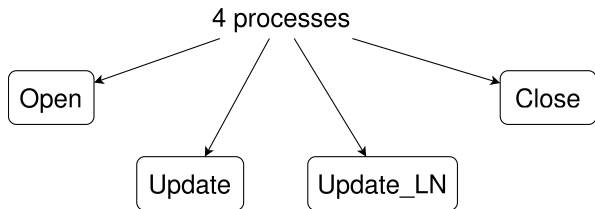


Modelling a transaction



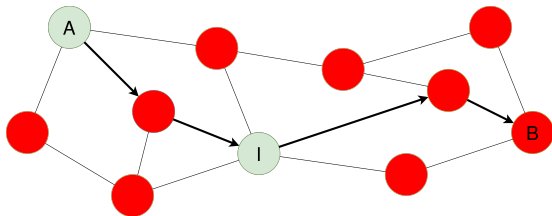
Transaction is represented by a quadruplet $\rightarrow tx = (pk_1, pk_2, b, h)$

Modelling a payment channel



Challenge: passing the state from one process to the other.

Modelling the whole network



Threat model:

- Honest agents communicate via authenticated and secret channels
- All agents can be compromised

Modelling the properties

- Indiana cannot lose money
 - **No money blocked:** Indiana can always close the channel
 - **No punishment possible:** Attacker cannot punish Indiana
 - **Defense against old states:** Indiana can punish old transactions
 - **Unforgeability:** attacker cannot forge transaction
- Indiana gets the fee
 - **Atomicity:** when Alice has paid, Indiana is able to debit

Difficulties

Obstacle encountered	Solution adopted
Time	Not modeling it
Passing the state	Using events
Liveness property	Tweak it into a correspondance property
Unbounded number of agents	Reduction to a bounded model

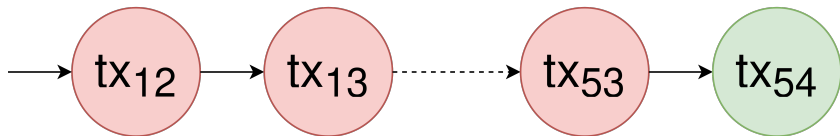
Trick



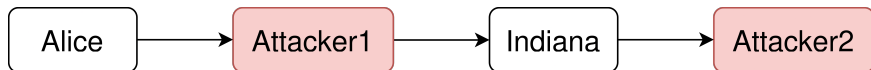
Liveness property: Indiana always holds a non-revoked transaction



Correspondance property: if transaction n Indiana holds is revoked, Indiana holds transaction $n+1$



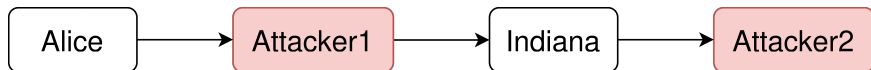
Reduced model



Theorem

Attack on the full network \implies attack on a 4-agent chain + oracles

Reduced model



Theorem

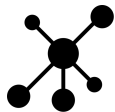
Attack on the full network \implies attack on a 4-agent chain + oracles

The attacker can simulate processes thanks to oracles.

```

let signing_oracle(sk: private_key)
  in(public, tx: transaction);
  event oracle_signs(tx, sk);
  out(public, sign(tx, sk)).
  
```

Conclusion



Modeling the LN protocol



Expressing all properties as correspondence property



Using a reduced model and a pen-and-paper proof



Next step: take time into account