Formal Analysis of Widevine DRM/EME

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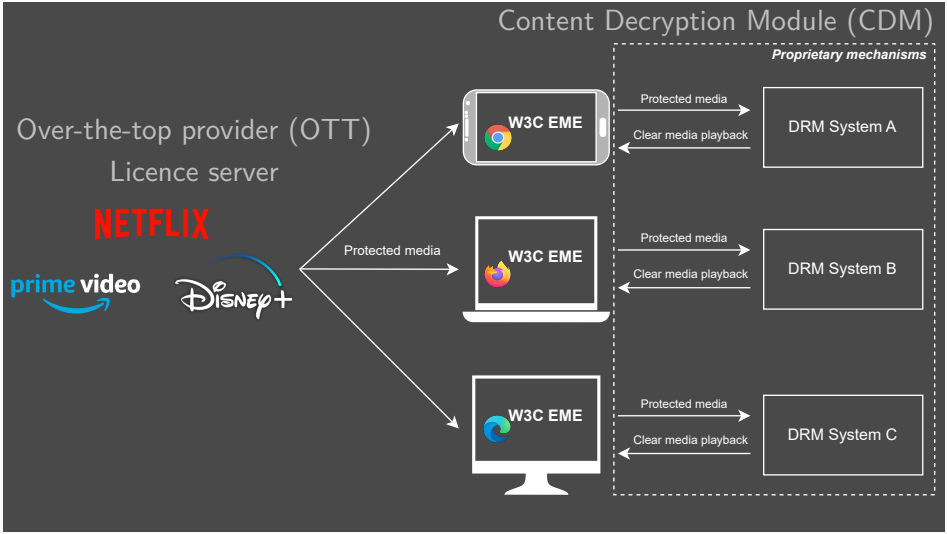
DRM Systems

- Digital Rights Management:
  restrict uses of digital content – prevent copy, *etc.*
- Used for music, books, video games, video streaming...

FairPlay

DRM Systems


Content Decryption Module (CDM)
Proprietary mechanisms
Over-the-top provider (OTT)
Licence server
NETFLIX
prime video
Disney+
Protected media
W3C EME
Protected media
Clear media playback
DRM System A
W3C EME
Protected media
Clear media playback
DRM System B
W3C EME
Protected media
Clear media playback
DRM System C

Encrypted Media Extension (EME)

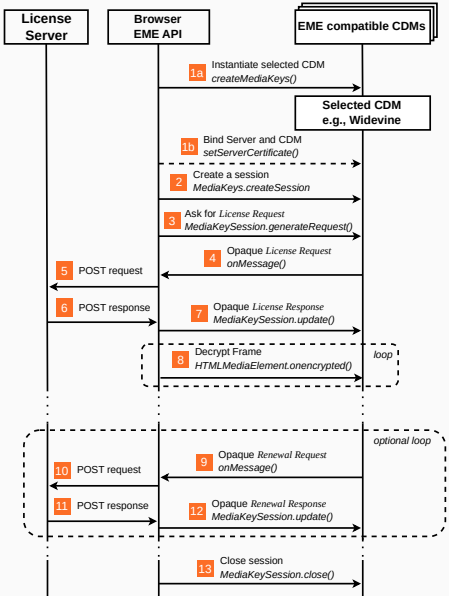
- EME: A standard defined by the World Wide Web Consortium (W3C)
- An API to make DRM use in browsers more uniform
- Integrated into all major browsers
- An "opaque" specification

Contribution

- Our goal: formally study the security of EME instantiated by Widevine
- Reverse engineer the proprietary Widevine protocol
- Define security properties (not present in EME spec)
- Model and prove in Tamarin

EME Specification

- Standard defines EME workflow and messages:
  - Initiate session
  - Initial Licence request/response
  - Licence Renewal request/response
  - Using a licence to decrypt
- Does not specify the content of messages or internal behaviour of the CDM
  $\rightarrow$ they are proprietary and implementation-specific



Reverse engineering of Widevine EME messages – Initial Licence

EME Message	Widevine EME Message Content
Licence Request	$\langle \text{reqID}, \text{nonce}, \{\text{clientID}\}_{\text{privacyK}}, \{\text{privacyK}\}_{\text{serviceCert}}, \text{keyID}, t_1 \rangle = \textit{req}$ + signature of *req* with deviceK
Licence Response	$\{\text{sessionK}\}_{\text{deviceK}}$, $\langle \text{reqID}, t_1, \Delta t, \text{keyID}, \{\text{contentK}\}_{\text{assetK}}, \{\text{nonce}\}_{\text{contentK}}, \text{policy} \rangle = \textit{resp}$ + MAC of *resp* with macK_S

- Hierarchy of keys:

$$\text{deviceK} \rightarrow \text{sessionK} \rightarrow \text{assetK}, \text{macK}_S, \text{macK}_C \rightarrow \text{contentK}$$
$$\text{assetK}, \text{macK}_S, \text{macK}_C = \mathit{KDF}(\mathit{tags}, \mathit{req}, \text{sessionK})$$

- nonce to ensure freshness
- Timestamps and time-to-live to control licence expiration

Reverse engineering of Widevine EME messages – Licence Renewal

EME Message	**Widevine EME Message Content**
Renewal Request	$\langle \text{reqID}, \{\text{clientID}\}_{\text{privacyK}'}, \{\text{privacyK}'\}_{\text{serviceCert}}, t_1, t_2, \text{ctr}, \text{nonce}' \rangle$ + MAC with macK_C
Renewal Response	$\langle \text{reqID}, t_1, t_2, \text{ctr}+1, \text{policy}', \Delta t, \{\text{nonce}'\}_{\text{contentK}} \rangle$ + MAC with macK_S

- Counter ctr to ensure correspondence between request/response
- Two slightly different versions:
  nonce′ in renewal is present on Android, but absent on desktop

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Security Goal 1: Key Confidentiality

Content decryption keys remain secret.

Security goals for initial licences

Security Goal 2: Integrity

The CDM must load initial licence responses *as they were generated* by the OTT.

Security Goal 3: Freshness

A given licence response can be loaded at most once,
and only by the CDM generating the corresponding request.

Security Goal 4: Expiration

In the initial phase, the CDM can use a decryption key at time t
only if the OTT granted a licence for it expiring at time $t_0 + \Delta t \geq t$.

Security goals for licence renewal

Security Goal 5: Integrity

The CDM must load renewal responses *as they were generated* by the OTT.

Security Goal 6: Freshness

A given renewal response can be loaded at most once, and only by the right CDM, and a CDM loads at most one response per *renewal event*.

Security Goal 7: Expiration

In the renewal phase, the CDM can use a decryption key at time t only if the OTT granted a (renewed) licence for it expiring at time $t_0 + \Delta t \geq t$.

Formal analysis of Widevine/EME

- We analyse Widevine/EME using the Tamarin prover (both Android and desktop version)
- A fairly complex protocol to model:
  - two roles (CDM & Licence Server), unbounded sessions 👍
  - a stateful API 👍
  - with loops
  - and counters
  - and timers...

Modelling of time in Tamarin

- Widevine messages contain timestamps, and Goals 4 and 7 explicitly mention time
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- No built-in support for time-based properties in Tamarin
- We propose our own encoding of time
  - Time as an integer %t
  - Each protocol rule receives it In(%t), has event GTime(%t)
  - It can appear in protocol messages and lemmas:

$$\text{GTime}(\%t_1)@i \;\&\; \text{State}(\%t_2, \ldots)@i \;\Rightarrow\; \%t_1 << \%t_2 \ldots$$

  - Attacker chooses the time each rule is executed, a restriction forces consistent choices:

$$\#i < \#j \;\&\; \text{GTime}(\%t_1)@i \;\&\; \text{GTime}(\%t_2)@j \;\Rightarrow\; \%t_1 << \%t_2$$

Conclusion & future work

- We reverse-engineered the Widevine DRM protocol
- We propose definitions for the security of Widevine as an EME instance
- We formally analyse the protocol in Tamarin

Future work

- Privacy properties
- Model for dishonest Licence server
- Other DRM systems

Questions?